# Multi-Objective ABC-NM Algorithm for Multi-Dimensional Combinatorial Optimization Problem 

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#### Abstract

This article addresses the problem of converting a single-objective combinatorial problem into a multi-objective one using the Pareto front approach. Although existing algorithms can identify the optimal solution in a multi-objective space, they fail to satisfy constraints while achieving optimal performance. To address this issue, we propose a multi-objective artificial bee colony optimization algorithm with a classical multi-objective theme called fitness sharing. This approach helps the convergence of the Pareto solution set towards a single optimal solution that satisfies multiple objectives. This article introduces multi-objective optimization with an example of a non-dominated sequencing technique and fitness sharing approach. The experimentation is carried out in MATLAB 2018a. In addition, we applied the proposed algorithm to two different real-time datasets, namely the knapsack problem and the nurse scheduling problem (NSP). The outcome of the proposed MBABCNM algorithm is evaluated using standard performance indicators such as average distance, number of reference solutions (NRS), overall count of attained solutions (TNS), and overall non-dominated generation volume (ONGV). The results show that it outperforms other algorithms.


Keywords: artificial bee colony; Nelder-Mead; multi-objective optimization; 0-1 knapsack problem; nurse scheduling problem

MSC: 68Q17; 68Q25; 68Q30; 68Q87

## 1. Introduction

Multi-objective optimization is the method of finding a single optimal result which has the potential to satisfy more than one objective for the given problem. There are three possible situations in multi-objective optimization problems [1,2]:

- Diminish all the objective functions;
- Increase all the objective functions;
- Diminish a few objectives and increase other objective functions.

$$
\begin{equation*}
\operatorname{Max} f(x)=\min (-f(x)) \tag{1}
\end{equation*}
$$

A multi-objective optimization problem (MOP) aims to determine a better compromising solution than a solitary individual. The vector of the decision variable $\overrightarrow{x^{*}} \in F$ is Pareto optimal when there is no other decision variable $\vec{x} \in F$ such that $f_{i}(\vec{x}) \leq f_{i}\left(\overrightarrow{x^{*}}\right)$, $i=1,2, \ldots, n$. The vector $\overrightarrow{x^{*}}$ is determined as Pareto optimal or globally non-dominated if no other solution in the set can dominate [3]. The set of solutions thus produced is said to be a Pareto-optimal set. The optimal set is specified as the Pareto-optimal front. From the multi-objective set, the user can select an ideal solution [4].

The dominance relation is used to associate two individuals. For example, a solution $u$ is said to lead to another solution $v$ if and only if $f_{i}(u) \leq f_{i}(v)$, for $i=1,2, \ldots, n$ and $f_{i}(u)<f_{i}(v)$ for at least one $i$, it is known as a globally non-dominated set. No other solution among the set can dominate it. The Pareto dominance for the solution is to dominate other solutions, and should not get worse in any given objectives with strict efficiency than one of them [5]. The Pareto dominance among two solutions $u$ and $v$ can possibly occur in any one of these cases:

- The solution $u$ dominates solution $v$, denoted as $u_{i} \prec v_{i}$;
- The solution $u$ is dominated by solution $v$, denoted as $v_{i} \prec u_{i}$;
- Both the solutions $u$ and $v$ are not dominated by each other, and they are said to be non-dominated. It is denoted as $\neg\left(u_{i} \prec v_{i}\right) \wedge \neg\left(v_{i} \prec u_{i}\right)$.
Recently, several meta-heuristic algorithms have been introduced to address the multidimensional combinatorial optimization problem [6]. Some of the famous techniques, namely genetic algorithm [7], differential evolution [8], particle swarm optimization [9], grey wolf optimization [10], and firefly algorithm [11], have been applied in various realtime applications, including optimum design for a centrifugal pump [12,13], Optimizing Magnification Ratio for the Flexible Hinge Displacement Amplifier [14], clustering [15], economic load dispatch [16], and job scheduling [17], to determine optimal solutions. However, the algorithms must be reinforced while applying them to multi-objective problems [18]. In this work, we utilized the ABC algorithm, which is more robust in mathematical analysis and provides more adequate solutions than all other algorithms. However, the ABC algorithm consumes ample computation time due to inefficient search direction while handling multi-objective problems. To eradicate these issues, we introduced the NelderMead technique with non-dominated sorting and fitness allotment methods to address the multi-objective concerns.

The main theme of this work is discussed below:

- A novel algorithm, MBABC-NM, is proposed to improve the exploitation of the artificial bee colony (ABC) technique. The algorithm incorporates a modified nondominated sorting and fitness-sharing approach to handle multi-dimensional problems efficiently.
- The proposed MBABC-NM algorithm is tested on two different real-time datasets: the knapsack problem and the nurse scheduling problem.
- The algorithm's performance is compared with other state-of-the-art algorithms, like genetic algorithm, cyber swarm optimization, and particle swarm optimization.
- The results of the experiments demonstrate that MBABC-NM outperforms the compared algorithms significantly. This result suggests that the proposed algorithm can effectively solve real-world optimization problems.
The rest of the paper is structured such that Section 2 discusses modified non-dominated sorting and fitness-sharing techniques over the multi-dimensional problem; Section 3 illustrates the detailed working process of the proposed MBABC-NM algorithm; and Section 4 presents the experimental setup for NSP and the $0-1$ knapsack problem. The empirical study and the discussion of the results are shown in Section 5, while Section 6 summarizes the work and its future directions.


## 2. Methodology

### 2.1. Modified Non-Dominated Sorting

In modified non-dominant sorting, the algorithm divides the population $L, 1 \leq L \leq N$ fronts in decreasing order of their dominance $F=\left\{F_{1}, F_{2}, \ldots, F_{L}\right\}$. Each solution in a front is non-dominated by the other. Each individual in $F_{l}$ is conquered by at least one individual in its preceding front $F_{j}$. Non-dominated arrangement aids in arranging the solutions sequentially based on the dominance, as mentioned in the above relation [19]. It improves the search capability of the multi-objective approach by introducing modified non-dominated solutions into the search space. The detailed narrative of the modified non-dominated arrangement is discussed in Algorithm 1. This algorithm is specified as one function involved in Algorithm 3.

```
Algorithm 1: Non-Dominated Sort (Z)
Input: Z
For each individual \(a \in Z\) do
Individuals dominated by \(a\)
    \(P_{a} \leftarrow \varnothing\)
    \(P_{b} \leftarrow \varnothing\)
Solutions which dominate \(a\)
        \(C_{a} \leftarrow 0\)
        For each solution \(b \in Z\) do
                if \((a \prec b)\) then
                Add the individuals \(b\) to the set of solutions dominated by \(a\)
                \(P_{a} \leftarrow P_{a} \cup\{b\}\)
                else if \((b \prec a)\) then
                Increment the domination counter \(a\)
                \(C_{a} \leftarrow C_{a}+1\)
                End if
    end for
if \(C_{a}=0\) then
Assign non-dominance rank as 1 for individual \(a\)
        \(a_{\text {rank }} \leftarrow 1\)
        \(L_{1} \leftarrow L_{1} \cup\{a\}\)
End if
end for
Initialize front counter
\(u \leftarrow 1\)
While \(L_{u} \neq \varnothing\) do
        Members of next front \(K\)
        \(K \leftarrow \varnothing\)
        For each solution \(a \in L_{u}\) do
            For each solution \(b \in P_{a}\) do
            Decrement the dominant counter of \(b\)
            \(C_{b} \leftarrow C_{b}-1\)
                        if \(C_{b}=0\) then
                            Assign rank for the individual \(b\)
                            \(b_{\text {rank }} \leftarrow u+1\)
                            \(K \leftarrow K \cup\{b\}\)
                        End if
            end for
        end for
    \(u \leftarrow u+1\)
    \(L_{u} \leftarrow K\)
The dominant solution of \(L_{u}\) are stored in \(L_{u}\)
```

This section discusses a modified non-domination sorting process that helps improve the multi-objective algorithm's search capability. In addition, the fitness-sharing function that aids the population in exploring diverse groups based on individual similarity is discussed in Section 2.2.

### 2.2. Fitness Sharing

Fitness sharing in evolutionary computing is used for isolating the population into diverse groups based on individual similarity [1]. It transforms an individual's fitness into the shared fitness value; usually, it is a lower value than the original. Only a limited amount of fitness value is available in each niche, and individuals in the same niche will share fitness value. The shared fitness $f_{\text {shared }}(i)$ of food particle $i$ with fitness $f i t_{i}$ can be measured by

$$
\begin{equation*}
f_{\text {shared }}(i)=\frac{f i t_{i}}{n_{i}} \tag{2}
\end{equation*}
$$

where $n_{i}$ is the niche total, which counts the number of food particles with fitness $f i t_{i}$ shared. The niche count can be calculated by summating the distribution function over the swarm.

$$
\begin{equation*}
n_{i}=\sum_{j=1}^{F P} \varphi\left(d_{i j}\right) \tag{3}
\end{equation*}
$$

where $F P$ denotes number of food particles and $\left(d_{i j}\right)$ is the distance between the food particles $i$ and $j$. The sharing function $\varphi$ computes the relationship between two food particles. The sharing function returns one if the food particles are identical and return 0 if the distance $\left(d_{i j}\right)$ is greater than a threshold of dissimilarity value. The distribution function can be represented as

$$
\varphi\left(d_{i j}\right)=f(x)=\left\{\begin{align*}
1-\left(\frac{d_{i j}}{\theta_{r}}\right)^{\varphi}, & d<\theta_{r}  \tag{4}\\
0, & \text { otherwise }
\end{align*}\right.
$$

where $\theta_{r}$ represents the sharing radius, which defines the size of the niche and threshold of dissimilarity. The food particles within this sharing radius are like each other and share their fitness. $\varphi$ is the constant which normalizes the shape of the distribution function. $d_{i j}$ is the distance between two food particles measured based on genotypic or phenotypic resemblance. The genotypic similarity is based on bit-string and is usually measured using Hamming distance. The phenotypic resemblance measures accurate parameters available in the search space using Euclidean distance.

$$
\begin{equation*}
d(a, b)^{\psi}=\left(\left[\left(p_{a}-p_{b}\right)^{2}+\left(q_{a}-q_{b}\right)^{2}\right]^{\frac{1}{2}}\right)^{\psi} \tag{5}
\end{equation*}
$$

The Euclidean distance $d(a, b)$ is the distance between the nodes $a$ and $b,\left(p_{a}, q_{a}\right)$ are the coordinates of the node $a$, and $\left(p_{b}, q_{b}\right)$ are the coordinates of the node $b$. The minimum transmission energy $T E_{s o l_{i}}$ contains the near-optimal solution. Fitness distribution based on phenotypic resemblance provides an improved outcome compared to distribution based on genotypic similarity [20-22].

In our algorithm, every individual finds a new solution. If a new solution dominates the original individual, it is entered into the external archive. If both do not dominate, then solutions are chosen randomly. When many non-dominated solutions exceed the archive size, our proposed algorithm uses a niching technique to truncate the crowded member and maintain uniform distribution among the archive members. Maintaining diversity among archive members is a complex task. Thus, in our proposed algorithm, we incorporated a fitness-sharing technique based on niche count, to ensure the diversity of the population.

The niching method maintains diversity and permits the algorithm to examine multiple summits in parallel. It also prevents the algorithm from being stuck in the local optima of search space, and can be viewed as the subspace of the population. For each niche in
our proposed algorithm, the fitness is finite and shared among the population. It is the process of optimizing the entire domain set. Fitness sharing transforms the raw fitness of a solution into a shared one. It helps to sustain diversity among the population, and thus our algorithm explores a better search space. The proposed fitness sharing technique is shown in Algorithm 2. Algorithm 3 invokes Algorithm 2 as a function while processing the execution.

```
Algorithm 2: Fitness Sharing \(\left(L_{u}\right)\)
    Number of solutions in Front counter \(L\)
        \(g \leftarrow\left|L_{u}\right|\)
For \(k \leftarrow 1\) to \(g\) do
        \(L_{u}\left(\right.\) Share \(\left._{k}\right) \leftarrow 0\)
        For each objective \(m\) do
        Sort population with respect to all objectives
            \(L_{u} \leftarrow \operatorname{sort}\left(L_{u}, m\right)\)
            \(L_{u}[1] \leftarrow \infty\)
            \(L_{u}[g] \leftarrow \infty\)
        For \(k \leftarrow 2\) to \(g-1\) do
        Calculate Shared fitness of the \(k^{\text {th }}\) solution with \(f_{i t}\)
            \(L_{u}\left(\right.\) Share \(\left._{k}\right) \leftarrow \frac{\text { fit }_{k}}{n_{k}}\)
                Niche count can be measured by
                \(n_{k} \leftarrow \sum_{j=1}^{|L|} \varphi\left(d_{k j}\right)\)
                The sharing function between two population elements can be measured using
                \(\varphi\left(d_{k j}\right) \leftarrow\left\{\begin{array}{r}1-\left(\frac{d_{k j}}{\theta_{r}}\right)^{\rho}, d<\theta_{r} \\ 0, \text { otherwise }\end{array}\right.\)
                End for
                End for
End for
```


## 3. Multi-Objective BABC-NM for a Multi-Dimensional Combinatorial Problem

Multi-Objective BABC-NM consists of an algorithm explained in this thesis and Algorithms 3-6 in this chapter. Algorithm 3 is modified based on a multi-objective perspective, and the pseudocode of the proposed MBABC-NM was described in detail in Algorithm 3. The working process of the formulated MBABC-NM is portrayed in Figure 1. The mapping process of this algorithm involves the following steps:

Initialization: The algorithm randomly generates a population of candidate solutions to the MDCOP. Each candidate solution is represented as a vector of decision variables.

Fitness Evaluation: The fitness of each candidate solution is evaluated by computing its objective function values. In multi-objective optimization, multiple objective functions usually need to be optimized simultaneously. Thus, the fitness of a candidate solution is represented as a vector of accurate function values.

Employed Bees: In this step, some bees are selected to perform the exploration process. The selected bees modify the solutions in the population by adding or subtracting a random value from the decision variables. It generates a new solution for each bee.

Onlooker Bees: Some other bees are selected to perform the exploitation process in this step. The selected bees choose solutions from the population based on their fitness and then modify them similarly to the employed bees. It generates a new solution for each onlooker bee.


Figure 1. Workflow of MBABC-NM.
Neighbourhood Mutation: In this step, the solutions generated by the employed and onlooker bees are subjected to a neighbourhood mutation process. It involves selecting a neighbourhood around each solution and generating a new solution within that neighbourhood.

Scout Bees: In this step, if a solution has not been improved after a certain number of iterations, it is considered a non-promising solution and is replaced by a new random solution generated by a scout bee.

Pareto Optimization: After generating the new solutions, the algorithm performs a Pareto optimization process to determine the best solutions. The Pareto optimization process identifies solutions not dominated by any other solution in the population.

Termination: The algorithm continues to iterate through steps 3 to 7 until a termination criterion is met. It could be a maximum number of iterations or a satisfactory level of solution quality.

```
Algorithm 3: MBABC-NM
Input
            FS: Number of Food Sources
            MI: Maximum iteration
                    Limit: number of predefined trials
Iter \(=0\)
Prepare the population
For \(i=1\) to FS do
        For \(j=1\) to \(S\) do
                            Produce \(x_{i, j}\) solution
                    \(x_{i, j} \leftarrow x_{m i n, j} \pm \operatorname{rand}(0,1) *\left(x_{m a x, j}-x_{m i n, j}\right)\)
                    Where \(x_{\text {min, } j}\) and \(x_{\text {max, }}\) are the min and max bound of the dimension \(j\)
                    \(\hat{x}_{i, j} \leftarrow \operatorname{BinaryConv}\left(x_{i, j}\right)\) using Algorithm 5
                    For \(h=1\) to \(M\) do
                            Evaluate the fitness of the population for a \(M\) number of Objectives
                        \(f_{h} \leftarrow f_{h}\left(\hat{x}_{i, j}\right)\)
                    End for
                        \(\operatorname{trial}(i) \leftarrow 0\)
```


## End for

## End for

```
iter \(\leftarrow 1\)
Repeat
\{
/ /*Employed Bee process*//
For each food source ido
Create new individual \(v_{i}\) using
\(v_{i, j} \leftarrow x_{i, j}+\varnothing_{i, j}\left(x_{i, j}-x_{k, j}\right)\)
\(\hat{v}_{i, j} \leftarrow \operatorname{BinaryConv}\left(v_{i, j}\right)\) using Algorithm 5
Evaluate \(f\left(\hat{v}_{i}\right)\)
Select between \(f\left(\hat{v}_{i}\right)\) and \(f\left(\hat{x}_{i}\right)\) using greedy method If \(f\left(\hat{v}_{i}\right)<f\left(\hat{x}_{i}\right)\)
\[
\begin{array}{r}
x_{i} \leftarrow v_{i} \\
f\left(\hat{x}_{i}\right) \leftarrow f\left(\hat{v}_{i}\right) \\
\operatorname{trial}(i) \leftarrow 0 \\
\operatorname{trial}(i) \leftarrow \operatorname{trial}(i)+1
\end{array}
\]
Else
End if
End For
```

/ /*Onlooker Bee Phase* / /
If iter $=1$
Set $r=0, i=1$;
While ( $\mathrm{r} \leq \mathrm{FS}$ )
Calculate Probabilities for onlooker bees using Algorithm 4
If rand $(0,1)<$ Pro $_{i}$

$$
r \leftarrow r+1
$$

For each food source, $i$ do
Generate new individual $v_{i}$ using Algorithm 6
NM method ( $v_{i}$ )
$\hat{v}_{i, j} \leftarrow$ BinaryConv $\left(v_{i, j}\right)$ using Algorithm 5
Evaluate $f\left(\hat{v}_{i}\right)$
Select between $f\left(\hat{v}_{i}\right)$ and $f\left(\hat{x}_{i}\right)$ using greedy method
If $f\left(\hat{v}_{i}\right)<f\left(\hat{x}_{i}\right)$

$$
\begin{gathered}
x_{i} \leftarrow v_{i} \\
f\left(\hat{x}_{i}\right) \leftarrow f\left(\hat{v}_{i}\right) \\
\operatorname{trail}(i) \leftarrow 0
\end{gathered}
$$

```
Algorithm 3: Cont.
            Else
                        \(\operatorname{trial}(i) \leftarrow \operatorname{trial}(i)+1\)
                    End if
                    End For
                    End if
                    \(i \leftarrow(i+1) \bmod F S\)
```


## End while

Else
For each food source, $i$ do
Generate new individual $v_{i}$ using Algorithm 6
NM method $\left(L_{u}\right)$

$$
u \in L_{u}
$$

Divide $\left\{L_{u}\right\}$ into $\left|L_{u}\right|$ equal chunks
$S_{u} \leftarrow \frac{\left\{L_{u}\right\}}{\left|L_{u}\right|}$
$\forall L_{u i}, i \in 1,2, \ldots,\left|L_{u}\right|$
$T_{x i} \leftarrow \operatorname{Rank}\left(L_{u i}, S_{u i}\right)$
$T_{x i} \leftarrow$ Delete least rank individual $\left(T_{x i}\right)$
$v_{i} \leftarrow \operatorname{celltomat}\left\{T_{x i}\right\}$
End For
End if
//*Scout Bee Phase* / /
$q=\{i: \operatorname{trial}(i)=\max ($ trial $)\}$
If $\operatorname{trial}(q)>$ limit
Abandon the food source $x_{i}$
$x_{q, j} \leftarrow x_{\text {min }, j} \pm \operatorname{rand}(0,1) *\left(x_{\text {max }, j}-x_{\text {min }, j}\right)$
$\hat{x}_{q, j} \leftarrow \operatorname{Binary} \operatorname{Conv}\left(x_{q, j}\right)$ using Algorithm 5
For $h=1$ to $M$ do
Evaluate the fitness of the population for a $M$ number
of Objectives

$$
f_{h} \leftarrow f_{h}\left(\hat{x}_{q}\right)
$$

## End for

$$
\operatorname{trial}(q) \leftarrow 0
$$

## End if

Add the new solution obtained to $Z_{i}$
Non-Dominated Sort ( $Z_{i}$ ) using Algorithm 1

$$
L \leftarrow Z_{i}
$$

Fitness Sharing ( $L$ ) using Algorithm 2//density estimation where $L$ denotes dense population around the individual $i$
Memorize the best solution obtained so far
iter $\leftarrow$ iter +1
\}
Until iter $=M I$
Output: Optimal value of the objective function

```
Algorithm 4: Probability Computation
For \(i=1\) to \(F S\), do
            Compute the probability \(P_{i j}\) for the individual \(v_{i, j}\)
            \(\operatorname{Pro}_{i} \leftarrow \frac{\text { fit }_{i}}{\sum_{j=1}^{f S} \text { fit }_{j}}\)
            fit \(_{i} \leftarrow \begin{cases}\frac{1}{1+f_{i}}, & f_{i} \geq 0 \\ 1+\operatorname{abs}\left(f_{i}\right), & f_{i}<0\end{cases}\)
End for
```

```
Algorithm 5: \(\operatorname{BinaryConv}\left(x_{i, j}\right)\)
For \(i=1\) to \(F S\) do
        For \(j=1\) to \(S\) do
            \(\operatorname{bit}\left(x_{i, j}\right)=\sin \left(2 \pi x_{i, j} \cos \left(2 \pi x_{i, j}\right)\right)\)
                        \(\hat{x}_{i, j}= \begin{cases}1, & \text { if bit }\left(x_{i, j}\right)>0 \\ 0, & \text { otherwise }\end{cases}\)
        End for
End for
```

```
Algorithm 6: NM method ( \(v_{i}\) )
Generate new food source \(v_{i}\) using modified NM technique
Let \(v_{i}\) denotes list of vertices
\(\mathrm{r}, \mu, \lambda\) and \(\zeta\) are the coefficients of reflection, expansion, contraction, and shrinkage
\(f\) is the objective function to be minimized
For \(i=1,2, \ldots, n+1\) vertices, do
    Arrange the values from lowest fit value \(f\left(v_{1}\right)\) to highest fit value \(f\left(v_{n+1}\right)\)
\(f\left(v_{1}\right) \leq f\left(v_{2}\right) \leq \ldots \leq f\left(v_{n+1}\right)\)
    Compute mean for best two summits
    \(v_{m} \leftarrow \sum \frac{v_{i}}{n}\), where \(i=1,2, \ldots, n\)
    \(/ / *\) Likeness point \(v_{r}^{*} / /\)
        \(v_{r} \leftarrow v_{m}+\mathrm{r}\left(v_{m}-v_{n+1}\right)\)
        If \(f\left(v_{1}\right) \leq f\left(v_{r}\right) \leq f\left(v_{n}\right)\) then
        \(v_{n} \leftarrow v_{r}\) and go to end condition
        End if
    \(/ / *\) Enlargement point \(v_{e}{ }^{*} / /\)
        If \(f\left(v_{r}\right) \leq f\left(v_{1}\right)\) then
        \(v_{e} \leftarrow v_{r}+\mu\left(v_{r}-v_{m}\right)\)
        End if
        If \(f\left(v_{e}\right)<f\left(v_{r}\right)\) then
        \(v_{n} \leftarrow v_{e}\) and go to end condition
        Else
        \(v_{n} \leftarrow v_{r}\) and go to end condition
        End if
    //*Reduction point \(v_{c}{ }^{*} / /\)
        If \(f\left(v_{n}\right) \leq f\left(v_{r}\right) \leq f\left(v_{n+1}\right)\) then
        Compute outside reduction
        \(v_{c} \leftarrow \lambda v_{r}+(1-\lambda) v_{m}\)
    End if
    If \(f\left(v_{r}\right) \geq f\left(v_{n+1}\right)\) then
    Compute inside reduction
    \(v_{c} \leftarrow \lambda v_{n+1}+(1-\lambda) v_{m}\).
    End if
    If \(f\left(v_{r}\right) \geq f\left(v_{n}\right)\) then
    Contraction is done among \(v_{m}\) and the best vertex among \(v_{r}\) and \(v_{n+1}\).
    End if
    If \(f\left(v_{c}\right)<f\left(v_{r}\right)\) then
        \(v_{n} \leftarrow v_{c}\) and go to end condition
        Else go to Shrinkage part
        End if
        If \(f\left(v_{c}\right) \geq f\left(v_{n+1}\right)\) then
        \(v_{n+1} \leftarrow v_{c}\) and go to end condition
        Else go to Shrinkage part
        End if
```

```
Algorithm 6: Cont.
    //*Shrinkage part *//
    Shrink towards the best solution with new vertices
    \(v_{i} \leftarrow \zeta v_{i}+v_{1}(1-\zeta)\), where \(i=2, \ldots, n+1\)
End condition
    Arrange and rename the newly constructed simplex's summits according to their fit
    values, then carry on with the reflection phase.
```


## 4. Experimental and Environment Setup

This section specifies the experimental structure of the proposed approach and other techniques. In addition, the projected outcomes with other methods are compared, to confirm the model's efficacy.

### 4.1. Experimental Setup

The proposed MBABC-NM algorithm to solve NSP and 0-1 knapsack problems, is demonstrated concisely in this division. The simulation is conducted on various optimization algorithms with similar environmental constraints, and the outcomes are analyzed. The technique proposed to handle NSP and 0-1 knapsack problems is implemented using a MATLAB 2018a tool under a Windows Intel I7 processor with 8GB of RAM. The experimental analysis will set the bounds of the formulated work. The parameters are considered based on the trial and error method. We used the standard dataset for both the NSP and the 0-1 knapsack problem. The compared algorithms are selected to ensure the performance of the formulated technique for the NSP in Table 1. The heuristic parameters and the consistent values are symbolized in Table 2.

Table 1. List of competitors' methods of comparing an NSP dataset for MBABC-NM.

| Type | Method | Reference |
| :---: | :---: | :---: |
| M1 | Multi-objective genetic algorithm: NSGA-II | Zhang et al., 2021 [23] |
| M2 | Multi-objective cyber swarm optimization algorithm | Yin et al., 2013 [24] |
| M3 | Multi-objective particle swarm optimization | Han et al., 2021 [25] |
| M4 | Multi-objective ABC | Li et al., 2015 [26] |

Table 2. Configuration parameters of MBABC-NM for experimental evaluation.

| Type | Method |
| :---: | :---: |
| number of bees | 100 |
| maximum iterations | 1000 |
| initialization technique | binary |
| stop criteria | maximum iterations |
| run | 20 |
| heuristic | Nelder-Mead method |
| likeness factor | $\alpha>0$ |
| enlargement factor | $\gamma>1$ |
| reduction factor | $0>\beta>1$ |
| shrinkage factor | $0<\delta<1$ |

### 4.2. Standard 0-1 Knapsack Problem Dataset

This work performs experiments on standard instances of the 0-1 knapsack problem from OR-Library to evaluate the performance of the proposed algorithm MBABC-NM. We used nine different instance classes to illustrate the outcomes of the proposed approach. Each problem suite is classified based on the number of knapsack constraints and object items used. A detailed description of the problem suite is discussed in Table 3. Table 3, column 2 describes the number of knapsacks constraints available in the corresponding problem suite. Column three represents the number of available object items within it, and
column four shows the known optimum solution provided by the standard OR-Library. The compared algorithms are selected to ensure the efficacy of the formulated model on the 0-1 knapsack problem, as shown in Table 4.

This section discussed the experimental setup and dataset description used for implementing the proposed algorithm and other compared techniques. Moreover, the performance of the proposed algorithm with other compared techniques is discussed in Section 5.

Table 3. The features of MKP datasets for MBABC-NM.

| Instance | No. of Objectives | No. of Items |
| :--- | :---: | :---: |
| kn250_2 | 2 | 250 |
| kn250_3 | 3 | 250 |
| kn250_4 | 4 | 250 |
| kn500_2 | 2 | 500 |
| kn500_3 | 3 | 500 |
| kn500_4 | 4 | 500 |
| kn750_2 | 2 | 750 |
| kn750_3 | 3 | 750 |
| kn750_4 | 4 | 750 |

Table 4. List of competitors' techniques to associate MKP dataset for MBABC-NM.

| Type | Method | Reference |
| :---: | :---: | :---: |
| M1 | Pareto evolutionary algorithm | Luo et al., 2019 [27] |
| M2 | GRASP | Yuan et al., 2021 [28] |
| M3 | Genetic Tabu search for MKP | Alharbi et al., 2018 [29] |
| M4 | ACO for MKP | Fidanova et al., 2020 [30] |

## 5. Experimental Result Analysis and Discussion

### 5.1. Standard NSP Dataset

The experimental outcomes achieved by the MBABC-NM algorithm on solving the standard NSP dataset are presented in Tables 5 and 6. The performance of the proposed algorithm is compared with existing multi-objective algorithms listed in Table 1 for M1, M2, M3, and M4. The value present in the table specifies the ONGV value attained via the consistent system. The multi-objectives of the NSP are the minimization of cost, maximizing nurse preferences, and minimizing the deviation between the number of nurses required and the least numeral of nurses for the shift; day shift followed by night shift is not permissible. To legalize the proposed algorithm, we utilized 15 test cases of various sizes with multiple issues. It is proven that projected MBABC-NM accomplished maximum ONGV values for a maximum number of instances. The experimentation has been carried out on four different algorithms with the same simulation parameters.

Table 5 reviews the comparison and assessment of ONGV performance indicators attained by our proposed technique MBABC-NM associated with other methods, as shown in Table 1.

On comparing the mean values of ONGV for the NSP dataset, our proposed MBABCNM outperforms existing algorithms for smaller datasets, with $14.99 \%$ against genetic NSGA, $59.67 \%$ against the cyber swarm, $28.70 \%$ against PSO, and $63.24 \%$ against the MABC algorithm. Our proposed MBABC-NM also outperforms existing algorithms for medium-sized datasets, with $84.75 \%$ against genetic NSGA, $23.12 \%$ against the cyber swarm, $60.43 \%$ against PSO, and $54.21 \%$ against the MABC algorithm. For larger-sized datasets, it achieved $43.15 \%$ against genetic NSGA, $44.27 \%$ against the cyber swarm, $25.14 \%$ against PSO, and $16.50 \%$ against the MABC algorithm.

Table 6 reviews the comparison and valuation of the SP performance indicator and shows the proposed MBABC-NM with another competitor's technique, as shown in

Table 1. Our proposed algorithm achieved minimized Euclidean distance among the Pareto solutions.

Table 5. Experimental result of NSP dataset in terms of ONGV.

| Case | Type | Instance | MBABC-NM | M1 | M2 | M3 | M4 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| C-1 | N25 | 1 | 135 | 121 | 92 | 104 | 53 |
| C-1 | N25 | 7 | 132 | 125 | 86 | 106 | 63 |
| C-1 | N25 | 12 | 131 | 119 | 91 | 115 | 51 |
| C-1 | N25 | 19 | 128 | 119 | 88 | 101 | 50 |
| C-1 | N25 | 25 | 128 | 122 | 83 | 104 | 56 |
| C-2 | N25 | 2 | 143 | 118 | 86 | 99 | 67 |
| C-2 | N25 | 5 | 142 | 121 | 80 | 113 | 69 |
| C-2 | N25 | 9 | 136 | 124 | 85 | 116 | 69 |
| C-2 | N25 | 15 | 149 | 124 | 78 | 115 | 60 |
| C-2 | N25 | 27 | 146 | 124 | 79 | 99 | 63 |
| C-3 | N25 | 1 | 145 | 123 | 77 | 97 | 61 |
| C-3 | N25 | 3 | 150 | 125 | 82 | 97 | 65 |
| C-3 | N25 | 16 | 151 | 125 | 77 | 99 | 71 |
| C-3 | N25 | 27 | 146 | 121 | 91 | 113 | 73 |
| C-3 | N25 | 35 | 151 | 117 | 93 | 107 | 70 |
| C-4 | N25 | 5 | 139 | 122 | 92 | 98 | 63 |
| C-4 | N25 | 10 | 136 | 117 | 88 | 112 | 71 |
| C-4 | N25 | 25 | 150 | 120 | 78 | 111 | 65 |
| C-4 | N25 | 38 | 151 | 121 | 97 | 110 | 59 |
| C-4 | N25 | 41 | 135 | 122 | 79 | 99 | 72 |
| C-5 | N25 | 7 | 150 | 122 | 78 | 97 | 59 |
| C-5 | N25 | 11 | 127 | 118 | 92 | 107 | 70 |
| C-5 | N25 | 30 | 135 | 120 | 80 | 114 | 61 |
| C-5 | N25 | 42 | 135 | 121 | 91 | 104 | 71 |
| C-5 | N25 | 47 | 148 | 118 | 83 | 100 | 64 |
| C-6 | N50 | 1 | 192 | 40 | 90 | 109 | 73 |
| C-6 | N50 | 4 | 229 | 47 | 91 | 107 | 60 |
| C-6 | N50 | 12 | 222 | 35 | 87 | 125 | 73 |
| C-6 | N50 | 26 | 244 | 47 | 96 | 114 | 66 |
| C-6 | N50 | 29 | 223 | 41 | 87 | 126 | 76 |
| C-7 | N50 | 3 | 242 | 36 | 96 | 65 | 57 |
| C-7 | N50 | 6 | 248 | 42 | 90 | 60 | 60 |
| C-7 | N50 | 12 | 246 | 34 | 87 | 67 | 66 |
| C-7 | N50 | 26 | 233 | 36 | 88 | 62 | 65 |
| C-7 | N50 | 36 | 214 | 39 | 89 | 72 | 61 |
| C-8 | N50 | 4 | 251 | 43 | 95 | 55 | 71 |
| C-8 | N50 | 9 | 255 | 48 | 98 | 74 | 55 |
| C-8 | N50 | 15 | 249 | 34 | 97 | 65 | 58 |
| C-8 | N50 | 40 | 196 | 37 | 87 | 57 | 57 |
| C-8 | N50 | 47 | 228 | 47 | 88 | 57 | 73 |
| C-9 | N60 | 5 | 225 | 36 | 94 | 63 | 61 |
| C-9 | N60 | 10 | 210 | 49 | 89 | 60 | 58 |
| C-9 | N60 | 23 | 207 | 33 | 99 | 73 | 63 |
| C-9 | N60 | 29 | 203 | 41 | 91 | 65 | 72 |
| C-9 | N60 | 40 | 183 | 37 | 100 | 73 | 67 |
| C-10 | N60 | 6 | 196 | 49 | 94 | 76 | 58 |
| C-10 | N60 | 14 | 180 | 47 | 90 | 65 | 66 |
| C-10 | N60 | 20 | 208 | 49 | 95 | 54 | 64 |
| C-10 | N60 | 32 | 184 | 42 | 91 | 64 | 60 |
| C-10 | N60 | 41 | 218 | 39 | 92 | 69 | 63 |
| C-11 | N60 | 2 | 349 | 82 | 137 | 123 | 129 |
| C-11 | N60 | 8 | 374 | 98 | 151 | 126 | 121 |
| C-11 | N60 | 14 | 316 | 83 | 144 | 113 | 111 |
| C-11 | N60 | 20 | 364 | 96 | 145 | 118 | 118 |
| C-11 | N60 | 32 | 292 | 96 | 139 | 112 | 134 |

Table 5. Cont.

| Case | Type | Instance | MBABC-NM | M1 | M2 | M3 | M4 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| C-12 | N60 | 3 | 327 | 98 | 140 | 121 | 115 |
| C-12 | N60 | 12 | 335 | 94 | 151 | 121 | 125 |
| C-12 | N60 | 19 | 351 | 98 | 145 | 120 | 124 |
| C-12 | N60 | 23 | 384 | 78 | 144 | 111 | 118 |
| C-12 | N60 | 34 | 289 | 98 | 140 | 121 | 107 |
| C-13 | N60 | 1 | 450 | 97 | 138 | 126 | 108 |
| C-13 | N60 | 4 | 438 | 87 | 141 | 118 | 109 |
| C-13 | N60 | 19 | 446 | 99 | 149 | 122 | 133 |
| C-13 | N60 | 29 | 347 | 81 | 152 | 126 | 109 |
| C-13 | N60 | 40 | 464 | 88 | 152 | 120 | 121 |
| C-14 | N60 | 5 | 335 | 100 | 141 | 108 | 121 |
| C-14 | N60 | 9 | 420 | 96 | 139 | 124 | 107 |
| C-14 | N60 | 15 | 400 | 90 | 146 | 116 | 115 |
| C-14 | N60 | 30 | 398 | 99 | 144 | 108 | 108 |
| C-14 | N60 | 43 | 483 | 94 | 140 | 117 | 110 |
| C-15 | N60 | 6 | 380 | 87 | 150 | 119 | 136 |
| C-15 | N60 | 15 | 433 | 87 | 141 | 123 | 125 |
| C-15 | N60 | 26 | 481 | 88 | 151 | 125 | 108 |
| C-15 | N60 | 35 | 477 | 90 | 151 | 124 | 136 |
| C-15 | N60 | 44 | 469 | 99 | 149 | 123 | 115 |

Table 6. Experimental results of NSP dataset in terms of SP.

| Case | Nurse | Instance | MBABC-NM | M1 | M2 | M3 | M4 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| C-1 | N25 | 1 | $1.1 \times 10^{-4}$ | $3.9 \times 10^{-4}$ | $2.0 \times 10^{-4}$ | $2.1 \times 10^{-4}$ | $4.2 \times 10^{-5}$ |
| C-1 | N25 | 7 | $1.1 \times 10^{-5}$ | $8.8 \times 10^{-4}$ | $2.4 \times 10^{-4}$ | $9.2 \times 10^{-5}$ | $3.1 \times 10^{-4}$ |
| C-1 | N25 | 12 | $1.1 \times 10^{-4}$ | $9.8 \times 10^{-4}$ | $2.2 \times 10^{-4}$ | $9.4 \times 10^{-5}$ | $8.4 \times 10^{-5}$ |
| C-1 | N25 | 19 | $1.9 \times 10^{-4}$ | $7.0 \times 10^{-4}$ | $1.1 \times 10^{-4}$ | $8.1 \times 10^{-6}$ | $1.3 \times 10^{-5}$ |
| C-1 | N25 | 25 | $1.1 \times 10^{-4}$ | $2.6 \times 10^{-4}$ | $1.0 \times 10^{-4}$ | $5.6 \times 10^{-4}$ | $3.3 \times 10^{-4}$ |
| C-2 | N25 | 2 | $1.8 \times 10^{-4}$ | $3.1 \times 10^{-4}$ | $4.7 \times 10^{-4}$ | $5.5 \times 10^{-4}$ | $2.3 \times 10^{-4}$ |
| C-2 | N25 | 5 | $9.7 \times 10^{-5}$ | $6.1 \times 10^{-4}$ | $3.4 \times 10^{-4}$ | $1.1 \times 10^{-4}$ | $2.9 \times 10^{-5}$ |
| C-2 | N25 | 9 | $7.4 \times 10^{-5}$ | $8.2 \times 10^{-4}$ | $4.0 \times 10^{-4}$ | $3.4 \times 10^{-4}$ | $9.5 \times 10^{-6}$ |
| C-2 | N25 | 15 | $6.6 \times 10^{-5}$ | $9.0 \times 10^{-5}$ | $3.5 \times 10^{-4}$ | $2.5 \times 10^{-4}$ | $3.3 \times 10^{-4}$ |
| C-2 | N25 | 27 | $3.7 \times 10^{-5}$ | $4.3 \times 10^{-5}$ | $2.7 \times 10^{-4}$ | $1.6 \times 10^{-4}$ | $1.0 \times 10^{-4}$ |
| C-3 | N25 | 1 | $7.2 \times 10^{-5}$ | $7.6 \times 10^{-4}$ | $4.4 \times 10^{-4}$ | $7.4 \times 10^{-5}$ | $1.0 \times 10^{-4}$ |
| C-3 | N25 | 3 | $1.4 \times 10^{-4}$ | $9.0 \times 10^{-4}$ | $1.0 \times 10^{-4}$ | $4.7 \times 10^{-4}$ | $3.5 \times 10^{-4}$ |
| C-3 | N25 | 16 | $1.5 \times 10^{-4}$ | $3.2 \times 10^{-4}$ | $2.0 \times 10^{-4}$ | $1.7 \times 10^{-4}$ | $1.4 \times 10^{-4}$ |
| C-3 | N25 | 27 | $6.6 \times 10^{-5}$ | $2.3 \times 10^{-4}$ | $3.5 \times 10^{-4}$ | $2.3 \times 10^{-4}$ | $2.7 \times 10^{-4}$ |
| C-3 | N25 | 35 | $4.5 \times 10^{-5}$ | $5.6 \times 10^{-4}$ | $3.8 \times 10^{-4}$ | $1.9 \times 10^{-4}$ | $2.7 \times 10^{-4}$ |
| C-4 | N25 | 5 | $1.9 \times 10^{-4}$ | $5.1 \times 10^{-4}$ | $7.0 \times 10^{-5}$ | $5.6 \times 10^{-4}$ | $1.9 \times 10^{-5}$ |
| C-4 | N25 | 10 | $9.4 \times 10^{-5}$ | $9.1 \times 10^{-4}$ | $8.4 \times 10^{-5}$ | $5.0 \times 10^{-4}$ | $1.0 \times 10^{-4}$ |
| C-4 | N25 | 25 | $1.6 \times 10^{-4}$ | $6.9 \times 10^{-4}$ | $6.3 \times 10^{-5}$ | $5.4 \times 10^{-4}$ | $2.0 \times 10^{-4}$ |
| C-4 | N25 | 38 | $3.2 \times 10^{-5}$ | $4.6 \times 10^{-4}$ | $1.0 \times 10^{-4}$ | $7.8 \times 10^{-5}$ | $8.1 \times 10^{-5}$ |
| C-4 | N25 | 41 | $1.4 \times 10^{-4}$ | $5.3 \times 10^{-4}$ | $1.9 \times 10^{-4}$ | $7.2 \times 10^{-5}$ | $1.3 \times 10^{-4}$ |
| C-5 | N25 | 7 | $1.8 \times 10^{-4}$ | $3.8 \times 10^{-5}$ | $2.4 \times 10^{-4}$ | $9.9 \times 10^{-5}$ | $1.2 \times 10^{-4}$ |
| C-5 | N25 | 11 | $3.5 \times 10^{-5}$ | $4.3 \times 10^{-4}$ | $4.2 \times 10^{-4}$ | $2.4 \times 10^{-4}$ | $1.2 \times 10^{-4}$ |
| C-5 | N25 | 30 | $1.2 \times 10^{-4}$ | $2.3 \times 10^{-4}$ | $4.1 \times 10^{-4}$ | $4.9 \times 10^{-4}$ | $2.0 \times 10^{-4}$ |
| C-5 | N25 | 42 | $2.1 \times 10^{-5}$ | $2.5 \times 10^{-4}$ | $4.6 \times 10^{-4}$ | $1.0 \times 10^{-4}$ | $2.2 \times 10^{-5}$ |
| C-5 | N25 | 47 | $1.8 \times 10^{-4}$ | $2.9 \times 10^{-4}$ | $6.2 \times 10^{-5}$ | $1.4 \times 10^{-4}$ | $1.0 \times 10^{-4}$ |
| C-6 | N50 | 1 | $1.2 \times 10^{-4}$ | $4.4 \times 10^{-4}$ | $1.7 \times 10^{-4}$ | $5.3 \times 10^{-4}$ | $2.3 \times 10^{-4}$ |
| C-6 | N50 | 4 | $1.9 \times 10^{-4}$ | $6.8 \times 10^{-4}$ | $6.3 \times 10^{-5}$ | $2.5 \times 10^{-5}$ | $3.2 \times 10^{-4}$ |
| C-6 | N50 | 12 | $1.4 \times 10^{-4}$ | $3.2 \times 10^{-5}$ | $2.5 \times 10^{-5}$ | $2.5 \times 10^{-4}$ | $1.5 \times 10^{-4}$ |
| C-6 | N50 | 26 | $9.4 \times 10^{-5}$ | $8.4 \times 10^{-4}$ | $1.1 \times 10^{-4}$ | $2.6 \times 10^{-4}$ | $3.3 \times 10^{-4}$ |
| C-6 | N50 | 29 | $3.4 \times 10^{-5}$ | $2.2 \times 10^{-4}$ | $1.8 \times 10^{-4}$ | $2.9 \times 10^{-4}$ | $3.4 \times 10^{-4}$ |

Table 6. Cont.

| Case | Nurse | Instance | MBABC-NM | M1 | M2 | M3 | M4 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| C-7 | N50 | 3 | $9.2 \times 10^{-5}$ | $3.5 \times 10^{-4}$ | $1.9 \times 10^{-4}$ | $2.4 \times 10^{-4}$ | $3.5 \times 10^{-4}$ |
| C-7 | N50 | 6 | $1.3 \times 10^{-4}$ | $4.5 \times 10^{-4}$ | $8.8 \times 10^{-5}$ | $2.5 \times 10^{-4}$ | $3.5 \times 10^{-4}$ |
| C-7 | N50 | 12 | $3.1 \times 10^{-5}$ | $4.0 \times 10^{-4}$ | $2.2 \times 10^{-4}$ | $8.0 \times 10^{-5}$ | $9.3 \times 10^{-5}$ |
| C-7 | N50 | 26 | $3.4 \times 10^{-5}$ | $1.8 \times 10^{-4}$ | $1.3 \times 10^{-4}$ | $4.3 \times 10^{-4}$ | $4.1 \times 10^{-5}$ |
| C-7 | N50 | 36 | $1.2 \times 10^{-4}$ | $2.1 \times 10^{-4}$ | $2.7 \times 10^{-4}$ | $3.7 \times 10^{-4}$ | $9.1 \times 10^{-5}$ |
| C-8 | N50 | 4 | $9.2 \times 10^{-5}$ | $1.3 \times 10^{-4}$ | $4.9 \times 10^{-4}$ | $3.5 \times 10^{-4}$ | $2.7 \times 10^{-4}$ |
| C-8 | N50 | 9 | $1.9 \times 10^{-4}$ | $1.2 \times 10^{-5}$ | $4.0 \times 10^{-5}$ | $8.6 \times 10^{-5}$ | $3.7 \times 10^{-5}$ |
| C-8 | N50 | 15 | $1.8 \times 10^{-4}$ | $1.2 \times 10^{-4}$ | $9.3 \times 10^{-5}$ | $4.3 \times 10^{-4}$ | $1.6 \times 10^{-4}$ |
| C-8 | N50 | 40 | $9.1 \times 10^{-5}$ | $5.7 \times 10^{-4}$ | $3.1 \times 10^{-4}$ | $2.1 \times 10^{-4}$ | $1.2 \times 10^{-4}$ |
| C-8 | N50 | 47 | $2.0 \times 10^{-4}$ | $8.6 \times 10^{-4}$ | $2.1 \times 10^{-4}$ | $2.3 \times 10^{-5}$ | $3.1 \times 10^{-4}$ |
| C-9 | N60 | 5 | $1.0 \times 10^{-4}$ | $9.6 \times 10^{-4}$ | $2.0 \times 10^{-4}$ | $3.5 \times 10^{-5}$ | $2.8 \times 10^{-4}$ |
| C-9 | N60 | 10 | $4.4 \times 10^{-5}$ | $6.2 \times 10^{-4}$ | $1.3 \times 10^{-4}$ | $4.5 \times 10^{-4}$ | $1.3 \times 10^{-4}$ |
| C-9 | N60 | 23 | $1.6 \times 10^{-4}$ | $5.5 \times 10^{-4}$ | $3.2 \times 10^{-4}$ | $3.6 \times 10^{-4}$ | $1.3 \times 10^{-5}$ |
| C-9 | N60 | 29 | $9.8 \times 10^{-5}$ | $6.6 \times 10^{-4}$ | $6.9 \times 10^{-5}$ | $2.3 \times 10^{-4}$ | $6.7 \times 10^{-5}$ |
| C-9 | N60 | 40 | $1.2 \times 10^{-4}$ | $2.0 \times 10^{-4}$ | $7.8 \times 10^{-5}$ | $2.6 \times 10^{-4}$ | $3.0 \times 10^{-4}$ |
| C-10 | N60 | 6 | $3.9 \times 10^{-5}$ | $2.0 \times 10^{-5}$ | $4.1 \times 10^{-4}$ | $5.4 \times 10^{-5}$ | $3.3 \times 10^{-4}$ |
| C-10 | N60 | 14 | $7.2 \times 10^{-5}$ | $2.3 \times 10^{-4}$ | $1.9 \times 10^{-4}$ | $1.4 \times 10^{-4}$ | $1.4 \times 10^{-4}$ |
| C-10 | N60 | 20 | $1.5 \times 10^{-4}$ | $8.8 \times 10^{-4}$ | $4.2 \times 10^{-5}$ | $2.1 \times 10^{-4}$ | $2.4 \times 10^{-4}$ |
| C-10 | N60 | 32 | $1.5 \times 10^{-4}$ | $8.4 \times 10^{-4}$ | $3.0 \times 10^{-4}$ | $7.0 \times 10^{-5}$ | $1.4 \times 10^{-4}$ |
| C-10 | N60 | 41 | $3.6 \times 10^{-5}$ | $7.0 \times 10^{-4}$ | $9.8 \times 10^{-5}$ | $2.4 \times 10^{-4}$ | $2.8 \times 10^{-4}$ |
| C-11 | N60 | 2 | $4.0 \times 10^{-5}$ | $9.4 \times 10^{-4}$ | $4.9 \times 10^{-4}$ | $4.1 \times 10^{-4}$ | $7.3 \times 10^{-5}$ |
| C-11 | N60 | 8 | $1.4 \times 10^{-5}$ | $4.5 \times 10^{-4}$ | $1.8 \times 10^{-4}$ | $3.8 \times 10^{-4}$ | $1.7 \times 10^{-4}$ |
| C-11 | N60 | 14 | $8.7 \times 10^{-5}$ | $2.1 \times 10^{-4}$ | $3.5 \times 10^{-4}$ | $2.7 \times 10^{-4}$ | $2.0 \times 10^{-4}$ |
| C-11 | N60 | 20 | $1.9 \times 10^{-4}$ | $4.3 \times 10^{-4}$ | $4.8 \times 10^{-4}$ | $2.3 \times 10^{-4}$ | $4.7 \times 10^{-5}$ |
| C-11 | N60 | 32 | $5.2 \times 10^{-5}$ | $2.3 \times 10^{-4}$ | $8.0 \times 10^{-5}$ | $5.4 \times 10^{-4}$ | $3.1 \times 10^{-4}$ |
| C-12 | N60 | 3 | $5.8 \times 10^{-5}$ | $1.1 \times 10^{-4}$ | $3.7 \times 10^{-4}$ | $4.5 \times 10^{-5}$ | $2.9 \times 10^{-4}$ |
| C-12 | N60 | 12 | $1.9 \times 10^{-4}$ | $9.1 \times 10^{-4}$ | $2.8 \times 10^{-4}$ | $2.0 \times 10^{-4}$ | $1.5 \times 10^{-4}$ |
| C-12 | N60 | 19 | $3.4 \times 10^{-5}$ | $7.5 \times 10^{-4}$ | $1.0 \times 10^{-4}$ | $3.3 \times 10^{-4}$ | $2.6 \times 10^{-4}$ |
| C-12 | N60 | 23 | $1.8 \times 10^{-4}$ | $3.4 \times 10^{-4}$ | $2.9 \times 10^{-4}$ | $9.3 \times 10^{-5}$ | $7.2 \times 10^{-5}$ |
| C-12 | N60 | 34 | $1.2 \times 10^{-4}$ | $5.2 \times 10^{-4}$ | $4.0 \times 10^{-4}$ | $5.3 \times 10^{-5}$ | $1.7 \times 10^{-4}$ |
| C-13 | N60 | 1 | $5.4 \times 10^{-5}$ | $9.8 \times 10^{-4}$ | $3.3 \times 10^{-4}$ | $3.6 \times 10^{-4}$ | $3.5 \times 10^{-4}$ |
| C-13 | N60 | 4 | $1.3 \times 10^{-4}$ | $6.7 \times 10^{-4}$ | $2.3 \times 10^{-4}$ | $3.5 \times 10^{-4}$ | $2.0 \times 10^{-4}$ |
| C-13 | N60 | 19 | $6.6 \times 10^{-6}$ | $7.5 \times 10^{-4}$ | $1.9 \times 10^{-4}$ | $1.4 \times 10^{-5}$ | $2.0 \times 10^{-4}$ |
| C-13 | N60 | 29 | $2.0 \times 10^{-4}$ | $4.9 \times 10^{-4}$ | $1.3 \times 10^{-4}$ | $4.3 \times 10^{-4}$ | $1.5 \times 10^{-4}$ |
| C-13 | N60 | 40 | $1.1 \times 10^{-4}$ | $5.8 \times 10^{-4}$ | $1.9 \times 10^{-4}$ | $1.3 \times 10^{-4}$ | $1.3 \times 10^{-5}$ |
| C-14 | N60 | 5 | $1.1 \times 10^{-4}$ | $7.5 \times 10^{-4}$ | $2.6 \times 10^{-4}$ | $5.0 \times 10^{-4}$ | $3.3 \times 10^{-4}$ |
| C-14 | N60 | 9 | $2.0 \times 10^{-4}$ | $3.9 \times 10^{-4}$ | $1.7 \times 10^{-4}$ | $3.0 \times 10^{-4}$ | $3.0 \times 10^{-4}$ |
| C-14 | N60 | 15 | $1.1 \times 10^{-4}$ | $1.9 \times 10^{-4}$ | $8.2 \times 10^{-5}$ | $1.3 \times 10^{-4}$ | $3.2 \times 10^{-4}$ |
| C-14 | N60 | 30 | $2.3 \times 10^{-4}$ | $6.5 \times 10^{-4}$ | $2.1 \times 10^{-4}$ | $1.1 \times 10^{-4}$ | $3.3 \times 10^{-4}$ |
| C-14 | N60 | 43 | $1.5 \times 10^{-4}$ | $8.1 \times 10^{-4}$ | $4.4 \times 10^{-4}$ | $5.4 \times 10^{-4}$ | $1.9 \times 10^{-4}$ |
| C-15 | N60 | 6 | $2.1 \times 10^{-6}$ | $5.7 \times 10^{-5}$ | $2.7 \times 10^{-4}$ | $5.5 \times 10^{-4}$ | $3.2 \times 10^{-4}$ |
| C-15 | N60 | 15 | $1.5 \times 10^{-4}$ | $3.5 \times 10^{-4}$ | $3.4 \times 10^{-4}$ | $1.4 \times 10^{-5}$ | $1.2 \times 10^{-4}$ |
| C-15 | N60 | 26 | $1.2 \times 10^{-4}$ | $7.7 \times 10^{-4}$ | $3.8 \times 10^{-4}$ | $1.6 \times 10^{-5}$ | $2.6 \times 10^{-4}$ |
| C-15 | N60 | 35 | $2.2 \times 10^{-5}$ | $1.3 \times 10^{-4}$ | $2.2 \times 10^{-4}$ | $8.1 \times 10^{-5}$ | $1.0 \times 10^{-4}$ |
| C-15 | N60 | 44 | $9.6 \times 10^{-5}$ | $8.4 \times 10^{-5}$ | $1.5 \times 10^{-4}$ | $1.6 \times 10^{-4}$ | $9.7 \times 10^{-7}$ |

Comparing the mean values of SP for the NSP dataset, our proposed MBABC-NM outperforms existing algorithms for smaller datasets, with $80.90 \%$ against genetic NSGA, $68.48 \%$ against the cyber swarm, $60.46 \%$ against PSO, and $43.86 \%$ against MABC algorithm. Our proposed MBABC-NM also outperforms existing algorithms for medium-sized datasets, with 79.55\% against Genetic NSGA, 59.14\% against the cyber swarm, $68.37 \%$ against PSO, and $57.84 \%$ against the MABC algorithm. For larger-sized datasets, it achieved $74.30 \%$ against Genetic NSGA, $53.72 \%$ against the cyber swarm, $63.22 \%$ against PSO, and $32.98 \%$ against the MABC algorithm.

### 5.2. Standard 0-1 Knapsack Problem

The experimental outcome achieved by the MBABC-NM algorithm on solving the standard 0-1 knapsack problem was presented in Tables 7 and 8. The performance of the proposed algorithm was compared with existing multi-objective meta-heuristic methods listed in Table 3. The values in the table specify the mean value attained for number of the reference solution and the total number of solutions using the corresponding algorithm. Our proposed MBABC-NM obtained better results for a maximum number of instances shown in [31].

Table 7. Experimental results of MKP dataset in terms of NRS and TNS.

| Instance | NRS |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | MBABC-NM | M1 | M2 | M3 | M4 | MBABC-NM | M1 | M2 | M3 | M4 |  |
| kn250_2 | 288.45 | 125.77 | 120.09 | 3.77 | 61.9 | 304.96 | 156.16 | 201.62 | 194.21 | 198 |  |
| kn250_3 | 549.19 | 191.68 | 183.63 | 0.69 | 92.2 | 663.75 | 207.34 | 376.82 | 1472.45 | 925 |  |
| kn250_4 | 742.71 | 215.63 | 207.90 | 1.28 | 104.6 | 791.22 | 254.71 | 617.90 | 3958.62 | 2288 |  |
| kn500_2 | 5019.36 | 1505.53 | 1500.14 | 1222.24 | 1361.2 | 5198.89 | 1543.83 | 1761.34 | 257.31 | 1009 |  |
| kn500_3 | 6751.66 | 2986.63 | 2980.58 | 1827.03 | 2403.8 | 6935.97 | 3032.85 | 3601.36 | 2368.15 | 2985 |  |
| kn500_4 | $17,156.62$ | 4282.56 | 4277.53 | 2258.97 | 3268.2 | $17,255.48$ | 5455.92 | 4930.68 | 5705.94 | 5318 |  |
| kn750_2 | $18,236.52$ | 4247.75 | 4240.96 | 3765.41 | 4003.2 | $20,515.01$ | 6087.35 | 4525.30 | 6362.16 | 5444 |  |
| kn750_3 | $33,682.95$ | 8035.34 | 8029.33 | 5336.95 | 6683.1 | $34,520.30$ | 9102.70 | 8297.50 | 7915.18 | 8106 |  |
| kn750_4 | $58,129.46$ | $11,307.37$ | $11,299.73$ | 6515.64 | 8907.7 | $60,293.70$ | $13,065.30$ | $11,648.42$ | 6976.39 | 9312 |  |

Table 8. Experimental results of MKP dataset in terms of the reference solution and Davg.

| Instance | \|R| | Davg |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | MBABC-NM | M1 | M2 | M3 | M4 |
| kn250_2 |  | $2.10 \times 10^{-4}$ | $9.70 \times 10^{-3}$ | $3.20 \times 10^{-3}$ | $1.48 \times 10^{-2}$ | $7.50 \times 10^{-3}$ |
| kn250_3 |  | $3.50 \times 10^{-4}$ | $1.50 \times 10^{-3}$ | $4.60 \times 10^{-3}$ | $2.02 \times 10^{-2}$ | $1.06 \times 10^{-2}$ |
| kn250_4 |  | $1.00 \times 10^{-4}$ | $3.14 \times 10^{-3}$ | $3.10 \times 10^{-3}$ | $3.24 \times 10^{-2}$ | $6.32 \times 10^{-3}$ |
| kn500_2 |  | $7.20 \times 10^{-4}$ | $1.78 \times 10^{-2}$ | $4.50 \times 10^{-3}$ | $1.61 \times 10^{-2}$ | $2.36 \times 10^{-2}$ |
| kn500_3 |  | $6.00 \times 10^{-4}$ | $1.41 \times 10^{-2}$ | $2.20 \times 10^{-3}$ | $3.24 \times 10^{-2}$ | $1.28 \times 10^{-2}$ |
| kn500_4 | 33374 | $2.50 \times 10^{-3}$ | $1.01 \times 10^{-2}$ | $3.60 \times 10^{-3}$ | $5.60 \times 10^{-2}$ | $3.00 \times 10^{-3}$ |
| kn750_2 | 34890 | $6.40 \times 10^{-3}$ | $2.46 \times 10^{-2}$ | $6.80 \times 10^{-3}$ | $3.17 \times 10^{-2}$ | $2.15 \times 10^{-2}$ |
| kn750_3 | 74504 | $9.50 \times 10^{-3}$ | $3.12 \times 10^{-2}$ | $7.80 \times 10^{-3}$ | $2.80 \times 10^{-2}$ | $9.80 \times 10^{-2}$ |
| kn750_4 | 105161 | $7.20 \times 10^{-3}$ | $2.93 \times 10^{-2}$ | $1.32 \times 10^{-2}$ | $3.18 \times 10^{-2}$ | $1.01 \times 10^{-2}$ |

The experimentation has been carried out on four different algorithms with the same simulation parameters. The outcome attained by the proposed technique MBABC-NM and another competitor algorithm is presented in Table 7. The values in the table represent the number of reference solutions obtained for corresponding algorithms. TNS symbolizes an overall count of attained solutions, and NRS defines the number of reference solutions for the instances. Table 8 describes the experimental work gained by our projected technique MBABC-NM and another competitor algorithm. |R| represents several Pareto optimal or reference sets obtained for our proposed algorithm. Davg denotes the average distance between the non-dominated individual and the reference set.

Compared with an existing algorithm, our proposed algorithm MBABC-NM generated a maximum number of reference solutions from the total solutions, as illustrated in Figures 2 and 3. The mean value of NRS for a smaller dataset with 250 objects, our proposed algorithm had achieved $38 \%$ more than other competitor algorithms. The mean value of TNS was $40 \%$ against the competitor algorithm. For a medium dataset with 500 objects, the NRS was $29 \%$, and the TNS was $21 \%$ against the competitor algorithm. For a larger dataset with 750 objects, our proposed algorithm MBABC-NM achieved $43 \%$ of NRS and $38 \%$ of TNS against the competitor algorithm.


Figure 2. Performance of MBABC-NM w.r.t NRS.


Figure 3. Performance of MBABC-NM with respect to TNS.
On comparing the mean values of Davg for the 0-1 knapsack problem dataset, as shown in Figure 4, our proposed MBABC-NM outperforms existing algorithms for smaller datasets with 250 objects, as it achieved $95.07 \%$ against Local search, $93.94 \%$ against GRASP, $98.87 \%$ against Genetic Tabu search, and $97.30 \%$ against the ACO algorithm. For a mediumsized dataset with 500 objects, our proposed algorithm achieved $90.90 \%$ against Local search, $62.91 \%$ against GRASP, $96.34 \%$ against Genetic Tabu search, and $90.30 \%$ against the ACO algorithm. For a larger dataset with 750 objects, we achieved $72.85 \%$ against Local search, $16.90 \%$ against GRASP, $74.75 \%$ against Genetic Tabu search, and $82.17 \%$ against the ACO algorithm.


Figure 4. Performance of MBABC-NM with respect to average distance.

## 6. Conclusions

This paper proposed multi-objective BABC-NM on the multi-dimensional combinatorial problem. The proposed multi-objective BABC-NM with fitness sharing and modified non-dominated sorting algorithm have been incorporated. The experimentation is carried out on a MATLAB 2018a. In addition, we consider the experimental setup to assess the outcome of the projected approach MBABC-NM. The practical results and discussions on the obtained effects prove the significance of the projected work. In all three experimental methodology stages, the projected algorithm MBABC-NM's enhanced outcome was outclassed by attaining precise and satisfactory outcome factors. The projected multi-objective binary ABC with Nelder-Mead (MBABC-NM) is outstripped for all the test cases when associating with other standard classical algorithms. These studies indeed confirmed the competence of the projected algorithm in all perceptions. The proposed algorithm could be extended to handle more complex real-time optimization problems, including scheduling and resource allocation. In addition, the algorithm could be further optimized to reduce its computational complexity and improve its scalability.

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